

Math/CS 467/667: Lecture 3

Given a quadrature rule

$$\text{quad}(f) = \sum_{k=0}^{n-1} w_k f(x_k) \quad \text{such that} \quad \int_{-1}^1 f(x) dx \approx \text{quad}(f) \quad (1)$$

is exact when f is a polynomial of degree less than or equal N , consider the approximation

$$\int_a^b f(x) dx \approx \text{comp}(f, a, b, m) \quad \text{where} \quad \text{comp}(f, a, b, m) = \sum_{i=0}^{m-1} \text{quad}(g_i)$$

is the composite quadrature formula on $[a, b]$ over m subintervals given by

$$g_j(x) = \frac{h}{2} f\left(\frac{xh}{2} + a + hj + \frac{h}{2}\right) \quad \text{and} \quad h = \frac{b-a}{m}.$$

Before beginning our analysis of the composite quadrature formula $\text{comp}(f, a, b, m)$, we first prove a lemma that results from the monotonicity of $\text{quad}(f)$ when the weights w_k are positive but also holds, in general, when they have mixed signs. Note that there are examples of naturally occurring Newton–Cotes quadrature formulas for which some of the weights w_k turn out to be negative. In the case when quad is given by Gaussian quadrature we have $N = 2n - 1$ and the weights are positive.

Lemma 2. *There is a constant $c \geq 2$ depending only on n and the w_k 's such that*

$$|f| \leq M \quad \text{implies} \quad |\text{quad}(f)| \leq cM$$

Proof. By the triangle inequality

$$|\text{quad}(f)| \leq \sum_{k=0}^{n-1} |w_k| |f(x_k)| \leq cM \quad \text{where} \quad c = \sum_{k=0}^{n-1} |w_k|.$$

In the case the $w_k \geq 0$ for all k we further have that

$$c = \sum_{k=0}^{n-1} |w_k| = \sum_{k=0}^{n-1} w_k = \text{quad}(1) = \int_{-1}^1 1 \cdot dx = 2.$$

Therefore, take $c = 2$ when all the weights are non-negative and note that $c > 2$ when some of the weights are negative. This finishes the proof of the lemma. ////

Note under the assumption the weights w_k are non-negative the approximation quad is, in fact, monotone as can be seen as follows: Suppose $f(x) \leq g(x)$ for all x , then

$$\text{quad}(f) = \sum_{k=0}^{n-1} w_k f(x_k) \leq \sum_{k=0}^{n-1} w_k g(x_k) = \text{quad}(g).$$

This, in particular, implies $|\text{quad}(f)| \leq \text{quad}(|f|)$, which means the approximation of the area under the absolute value of a function is always larger than the absolute value of the approximation of its integral.

We characterize now the error in the composite quadrature formula by proving

Theorem 3. *If f has $N + 1$ continuous derivatives on the interval $[a, b]$ then the error*

$$E_m = \left| \int_a^b f(x)dx - \text{comp}(f, a, b, m) \right| = \mathcal{O}(h^{N+1}) \quad \text{as} \quad h \rightarrow 0.$$

Proof. Let $t_j = a + hj$ and note that

$$\int_a^b f(x)dx = \sum_{j=0}^{m-1} \int_{t_j}^{t_{j+1}} f(t)dt.$$

For each of the integrals over the intervals $[t_j, t_{j+1}]$ of length h appearing on the right hand side make the change of variables

$$t = \frac{-t_j(x-1)}{2} + \frac{t_{j+1}(x+1)}{2} = \frac{xh}{2} + a + hj + \frac{h}{2} \quad \text{and} \quad dt = \frac{h}{2}dx$$

to obtain

$$\int_{t_j}^{t_{j+1}} f(t)dt = \frac{h}{2} \int_{-1}^1 f\left(\frac{xh}{2} + a + hj + \frac{h}{2}\right)dx = \int_{-1}^1 g_j(x)dx.$$

We now use the fact that quad is exact for polynomials of degree less or equal N to obtain bounds on the error. By the triangle inequality

$$E_m \leq \sum_{j=0}^{m-1} \left| \int_{-1}^1 g_j(x)dx - \text{quad}(g_j) \right|. \quad (4)$$

Since f has $N + 1$ continuous derivatives and the maximum of the continuous function $f^{(N+1)}(x)$ is guaranteed to exist on the closed interval $[a, b]$, then we may define

$$M = \max \{ |f^{(N+1)}(x)| : x \in [a, b] \}.$$

Upon noting that g_j also has $N + 1$ continuous derivatives, it follows from Taylor's theorem that $g_j(x) = T_j(x) + R_j(x)$ where T_j is the Taylor polynomial of degree N expanded about $x = 0$ and R_j is the remainder given by

$$R_j(x) = \frac{x^{N+1}}{(N+1)!} g_j^{(N+1)}(\xi_j) \quad \text{for some } \xi_j \text{ between } 0 \text{ and } x.$$

Since $x \in [-1, 1]$ then $\xi_j \in [-1, 1]$. By the chain rule we obtain

$$\begin{aligned} |g_j^{(N+1)}(\xi_j)| &= \frac{h}{2} \left| \left(\frac{d}{dx} \right)^{N+1} f\left(\frac{xh}{2} + a + hj + \frac{h}{2}\right) \right|_{x=\xi_j} \\ &= \left(\frac{h}{2} \right)^{N+2} \left| f^{(N+1)}\left(\frac{\xi_j h}{2} + a + hj + \frac{h}{2}\right) \right| \\ &\leq \left(\frac{h}{2} \right)^{N+2} \max \{ |f^{(N+1)}(t)| : t \in [t_j, t_{j+1}] \} \leq \left(\frac{h}{2} \right)^{N+2} M. \end{aligned}$$

Consequently,

$$|R_j(x)| \leq \frac{|x|^{N+1}}{(N+1)!} \left(\frac{h}{2}\right)^{N+2} M \leq h^{N+2} B \quad \text{where} \quad B = \frac{1}{(N+1)!} \cdot \frac{M}{2^{N+2}}.$$

Plugging the Taylor polynomial and remainder into (4) and using the fact that **quad** is exact for polynomials of degree less than or equal N we obtain

$$E_m \leq \sum_{j=0}^{m-1} \left| \int_{-1}^1 R_j(x) dx \right| + \sum_{j=0}^{m-1} |\text{quad}(R_j)|.$$

At this point we use the monotonicity of the integral—the fact that the area under the absolute value of a curve is greater than the original area—to estimate

$$\left| \int_{-1}^1 R_j(x) dx \right| \leq \int_{-1}^1 |R_j(x)| dx \leq \int_{-1}^1 h^{N+2} B = 2h^{N+2} B.$$

Combining the above estimate with the Lemma 2 applied to $|\text{quad}(R_j)|$ yields

$$\begin{aligned} E_m &\leq \sum_{i=0}^{m-1} (2+c)h^{N+2} B = (2+c)mh^{N+2} B \\ &= (2+c)B(b-a)h^{N+1} = \mathcal{O}(h^{N+1}) \quad \text{as} \quad h \rightarrow 0. \end{aligned}$$

This finishes the proof of the theorem. ////

We remark in the case of Gaussian quadrature where $N = 2n - 1$ that the results of Theorem 3 may be simplified to obtain

$$E_m \leq 4B(b-a) \left(\frac{b-a}{m}\right)^{2n}.$$

In applications one typically chooses n fixed and then increases m until the desired error goals are met. It is, of course, possible to increase n as well. However, in doing so, one must remember that B also depends on n through M .

If f is an analytic function such that its Taylor series converges on a closed disk in the complex plane of radius r at every point $x \in [a, b]$, this means for complex ω that

$$\max \left\{ \frac{r^{2n}}{(2n)!} |f^{(2n)}(\omega)| : |\omega - x| \leq r \right\} \leq \max \left\{ |f(\omega)| : |\omega - x| = r \right\}$$

for all $n \geq 0$. It follows that $(2+c)B(2r)^{2n}(b-a)$ is bounded, say by A , and consequently it holds for every $h < 2r$ that

$$E_m \leq A \left(\frac{h}{2r}\right)^{2n} \rightarrow 0 \quad \text{exponentially as} \quad n \rightarrow \infty.$$

Thus, provided h is small enough, it is also possible—though less common—to meet any error bounds with exponential efficiency by taking n sufficiently large.